

Low-discrepancy constructions in the triangle

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ICERM 29th October 2014



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Introduction

- Problem : Numerical Integration over triangular domain using quasi-Monte Carlo (QMC) sampling.
- QMC in $[0, 1]^d$.

$$\mu = \int_{[0,1]^d} f(x) dx \quad \hat{\mu}_n = \frac{1}{n} \sum_{i=1}^n f(x_i)$$

- Koksma-Hlawka inequality

$$|\hat{\mu}_n - \mu| \leq D_n^*(x_1, \dots, x_n) \times V_{HK}(f)$$

- Recent work relating to general spaces by Aistleitner et al., (2012), Brandolini et. al (2013)



Motivation

- Need in computer graphics, genetic experimental studies, etc.
- Mapping by special functions/transformation from $[0, 1]^d$
Pillards and Cools (2005).
- Several notions of discrepancy on the triangle/simplex but no explicit constructions. Pillards and Cools (2005), Brandolini et. al (2013).



General Notions of Discrepancy

- The signed discrepancy of \mathcal{P} at the measurable set $S \subseteq \Omega \subset \mathbb{R}^d$ is

$$\delta_N(S; \mathcal{P}, \Omega) = \mathbf{vol}(S \cap \Omega) / \mathbf{vol}(\Omega) - A_N(S; \mathcal{P}) / N.$$

- The absolute discrepancy of points \mathcal{P} for a class \mathcal{S} of measurable subsets of Ω is

$$D_N(\mathcal{S}; \mathcal{P}, \Omega) = \sup_{S \in \mathcal{S}} D_N(S; \mathcal{P}, \Omega),$$

where

$$D_N(S; \mathcal{P}, \Omega) = |\delta_N(S; \mathcal{P}, \Omega)|.$$

- Standard QMC works with $\Omega = [0, 1)^d$ and takes for \mathcal{S} the set of anchored boxes $[0, \mathbf{a})$ with $\mathbf{a} \in [0, 1)^d$.



Discrepancy due to Brandolini et al. (2013)

- $\mathcal{S}_C = \{\mathcal{T}_{a,b,C} \mid 0 < a < \|A - C\|, 0 < b < \|B - C\|\}$
- The parallelogram discrepancy of points \mathcal{P} for $\Omega = \Delta(A, B, C)$ is

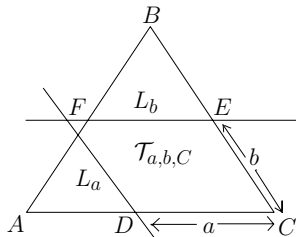
$$D_N^{\mathcal{P}}(\mathcal{P}; \Omega) = D_N(\mathcal{S}_{\mathcal{P}}; \mathcal{P}, \Omega)$$

for

$$\mathcal{S}_{\mathcal{P}} = \mathcal{S}_A \cup \mathcal{S}_B \cup \mathcal{S}_C.$$

Figure: The construction of the parallelogram

$$\mathcal{T}_{a,b,C} = CDFE$$





Discrepancy due to Pillards and Cools (2005)

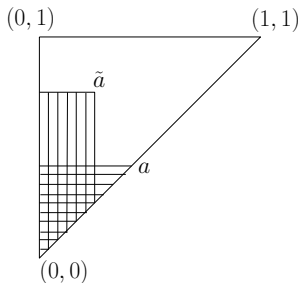
- $\Omega = \Delta((0,0)^T, (0,1)^T, (1,1)^T)$
- Their discrepancy

$$D_N^{PC}(\mathcal{P}; \Omega) = D_N(\mathcal{S}_I, \mathcal{P}, \Omega)$$

where

$$\mathcal{S}_I = \{[0, \mathbf{a}] \mid \mathbf{a} \in [0, 1)^2\}.$$

Figure: Star Discrepancy on the Simplex





Relationship between the discrepancies

Lemma 1

Let T_{PC} be the triangle from Pillards and Cools and for $N \geq 1$, let \mathcal{P} be the list of points $\mathbf{x}_1, \dots, \mathbf{x}_N \in T_{PC}$. Then

$$D_N^{PC}(\mathcal{P}, T_{PC}) \leq 2D_N^P(\mathcal{P}, T_{PC})$$

Proof

- $[0, a_1) \times [0, a_2) = [0, a_1) \times [0, 1) - [0, a_1) \times [a_2, 1)$
- Taking C to be the vertex $(0, 1)^T$ of T_{PC} ,
- $D_N^{PC}(\mathcal{P}; T_{PC}) \leq 2D_N(\mathcal{S}_C, \mathcal{P}, T_{PC}) \leq 2D_N^P(\mathcal{P}, T_{PC})$.



Triangular van der Corput construction

- van der Corput sampling of $[0, 1]$ the integer $n = \sum_{k \geq 1} d_k b^{k-1}$ in base $b \geq 2$ is mapped to $x_n = \sum_{k \geq 1} d_k b^{-k}$.
- Points $x_1, \dots, x_n \in [0, 1)$ have a discrepancy of $O(\log(n)/n)$.
- Our situation : 4-ary expansion.



Triangular van der Corput construction

- $n \geq 0$ in a base 4 representation $n = \sum_{k \geq 1} d_k 4^{k-1}$ where $d_k \in \{0, 1, 2, 3\}$

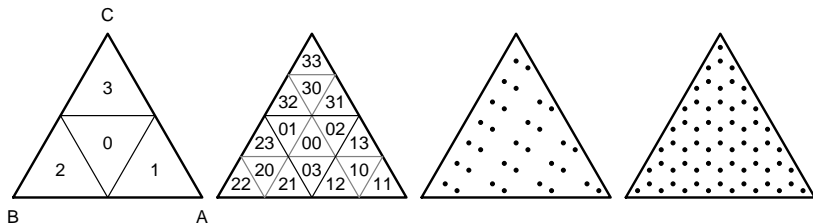


Figure: A labeled subdivision of $\Delta(A, B, C)$ into 4 and then 16 congruent subtriangles. Next are the first 32 triangular van der Corput points followed by the first 64. The integer labels come from the base 4 expansion.



Triangular van der Corput construction

- Computation : $T = \Delta(A, B, C)$

$$T(d) = \begin{cases} \Delta\left(\frac{B+C}{2}, \frac{A+C}{2}, \frac{A+B}{2}\right), & d = 0 \\ \Delta\left(A, \frac{A+B}{2}, \frac{A+C}{2}\right), & d = 1 \\ \Delta\left(\frac{B+A}{2}, B, \frac{B+C}{2}\right), & d = 2 \\ \Delta\left(\frac{C+A}{2}, \frac{C+B}{2}, C\right), & d = 3. \end{cases}$$

- This construction defines an infinite sequence of $f_T(i) \in T$ for integers $i \geq 0$.
- For an n point rule, take $\mathbf{x}_i = f_T(i-1)$ for $i = 1, \dots, n$.



Discrepancy Results

Theorem 1

For an integer $k \geq 0$ and non-degenerate triangle $\Omega = \Delta(A, B, C)$, let \mathcal{P} consist of $\mathbf{x}_i = f_{\Omega}(i-1)$ for $i = 1, \dots, N = 4^k$. Then

$$D_N^P(\mathcal{P}; \Omega) = \begin{cases} \frac{7}{9}, & N = 1 \\ \frac{2}{3\sqrt{N}} - \frac{1}{9N}, & \text{else.} \end{cases}$$



Discrepancy Results

Theorem 2

Let Ω be a nondegenerate triangle, and let \mathcal{P} contain points $\mathbf{x}_i = f_{\Omega}(s + i - 1)$, $i = 1, \dots, N = 4^k$, for a starting integer $s \geq 1$ and an integer $k \geq 0$. Then

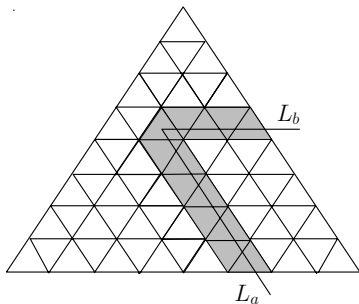
$$D_N^{\mathcal{P}}(\mathcal{P}; \Omega) \leq \frac{2}{\sqrt{N}} - \frac{1}{N}.$$



Proof of Theorem 2

Proof

- $\delta_N(S) = \sum_{j=0}^m \delta_N(S_j)$ where m is the number of subtriangles touching a boundary line of $\mathcal{T}_{a,b,c}$.
- $-1/N \leq \delta_N(S_j) \leq 1/N$.
- $D_N(S; \mathcal{P}) \leq m/N$
- $m \leq 2\sqrt{N} - 1$
- $D_N(S_C; \mathcal{P}) \leq (2\sqrt{N} - 1)/N$





Discrepancy Results

Theorem 3

Let Ω be a non-degenerate triangle and, for integer $N \geq 1$, let $\mathcal{P} = (\mathbf{x}_1, \dots, \mathbf{x}_N)$, where $\mathbf{x}_i = f_{\Omega}(i-1)$. Then

$$D_N^{\mathcal{P}}(\mathcal{P}; \Omega) \leq 12/\sqrt{N}.$$

Proof:

- Let $N = \sum_{j=0}^k a_j 4^j$ for some k , with $a_k \neq 0$.
- Let \mathcal{P}_j^l denote a set of 4^j consecutive points from \mathcal{P} , for $l = 1, \dots, a_j$ and $j \leq k$. These \mathcal{P}_j^l can be chosen to partition the N points \mathbf{x}_i . Fix any $S \in \mathcal{S}_{\mathcal{P}}$.



Proof of Theorem 3

- Now,

$$\delta_N(S; \mathcal{P}) = \frac{1}{N} \sum_{j=0}^k \sum_{l=1}^{a_j} 4^j \delta(S; \mathcal{P}'_j).$$

- Therefore from Theorem 2,

$$\begin{aligned} D_N(S; \mathcal{P}) = |\delta_N(S; \mathcal{P})| &\leq \frac{1}{N} \sum_{j=0}^k \sum_{l=1}^{a_j} 4^j \left(\frac{2}{2^j} - \frac{1}{4^j} \right) \leq \frac{1}{N} \sum_{j=0}^k a_j (2^{j+1} - 1) \\ &\leq \frac{3}{N} (2(2^{k+1} - 1) - (k + 1)) \leq \frac{12 \times 2^k}{N} \end{aligned}$$

and then $k \leq \log_4(N)$, gives $D_N(S; \mathcal{P}) \leq 12/\sqrt{N}$.

- Taking the supremum over $S \in \mathcal{S}_P$ yields the result.



Triangular Kronecker Lattice

- We use Theorem 1 of Chen and Travaglini (2007)
- This construction yields parallel discrepancy of $O(\log N/N)$

Definition 1

A real number θ is said to be *badly approximable* if there exists a constant $c > 0$ such that $n||n\theta|| > c$ for every natural number $n \in \mathbb{N}$ and $||\cdot||$ denotes the distance from the nearest integer.

Definition 2

Let a, b, c and d be integers with $b \neq 0, d \neq 0$ and $c > 0$, where c is not a perfect square. Then $\theta = (a + b\sqrt{c})/d$ is a *quadratic irrational number*.



Triangular Kronecker Lattice

- Let $\Theta = \{\theta_1, \dots, \theta_k\}$ be a set of $k \geq 1$ angles in $[0, 2\pi)$.
- Then let $\mathcal{A}(\Theta)$ be the set of convex polygonal subsets of $[0, 1]^2$ whose sides make an angle of θ_i with respect to the horizontal axis.

Theorem 1 (Chen and Travaglini (2007))

There exists a constant $C_\Theta < \infty$ such that for any integer $N > 1$ there exists a list $\mathcal{P} = (\mathbf{x}_1, \dots, \mathbf{x}_N)$ of points in $[0, 1]^2$ with

$$D_N(\mathcal{A}(\Theta); \mathcal{P}, [0, 1]^2) < C_\Theta \log(N)/N.$$



Triangular Kronecker Lattice

Lemma 2 (Davenport)

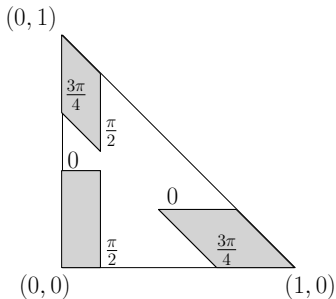
Suppose that the angles $\theta_1, \dots, \theta_k \in [0, 2\pi)$ are fixed. Then there exists $\alpha \in [0, 2\pi)$ such that $\tan(\alpha), \tan(\alpha - \pi/2), \tan(\alpha - \theta_1), \dots, \tan(\alpha - \theta_k)$ are all finite and badly approximable.



Triangular Kronecker Lattice

- $R = \Delta((0,0)^T, (0,1)^T, (1,0)^T)$.
- $\Theta = \{0, \pi/2, 3\pi/4\}$

Figure: Set of Angles for Kronecker Construction





Triangular Kronecker Lattice

Lemma 3

Let α be an angle for which $\tan(\alpha)$ is a quadratic irrational number. Then $\tan(\alpha)$, $\tan(\alpha - \pi/2)$ and $\tan(\alpha - 3\pi/4)$ are all finite and badly approximable.

- $\tan(3\pi/8) = 1 + \sqrt{2}$.
- $\tan(5\pi/8) = -1 - \sqrt{2}$.



Triangular Kronecker Lattice

Theorem 4

Let $N > 1$ be an integer and let R defined above be the triangle. Let $\alpha \in (0, 2\pi)$ be an angle for which $\tan(\alpha)$ is a quadratic irrational. Let \mathcal{P}_1 be the points of the lattice $(2N)^{-1/2}\mathbb{Z}^2$ rotated anticlockwise by angle α . Let \mathcal{P}_2 be the points of \mathcal{P}_1 that lie in R . If \mathcal{P}_2 has more than N points, let \mathcal{P}_3 be any N points from \mathcal{P}_2 , or if \mathcal{P}_2 has fewer than N points, let \mathcal{P}_3 be a list of N points in R including all those of \mathcal{P}_2 . Then there is a constant C with

$$D^P(\mathcal{P}_3; R) < C \log(N)/N.$$



Triangular Kronecker Lattice

Triangular Lattice Points

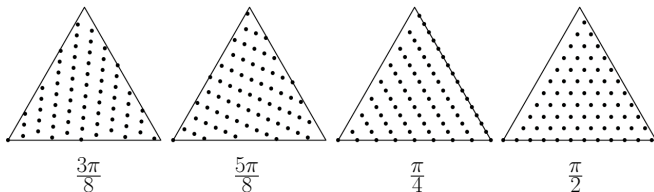


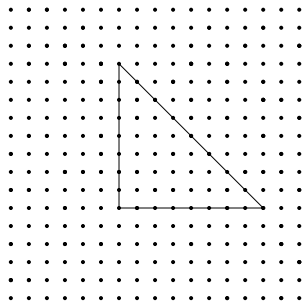
Figure: Triangular lattice points for target $N = 64$. Domain is an equilateral triangle. Angles $\frac{3\pi}{8}$ and $\frac{5\pi}{8}$ have badly approximable tangents. Angles $\frac{\pi}{4}$ and $\frac{\pi}{2}$ have integer and infinite tangents respectively and do not satisfy the conditions for discrepancy $O(\log(N)/N)$.



Construction Algorithm

Given a target sample size N , an angle α such as $3\pi/8$ satisfying Lemma 3. and a target triangle $\Delta(A, B, C)$,

- Take integer grid \mathbb{Z}^2

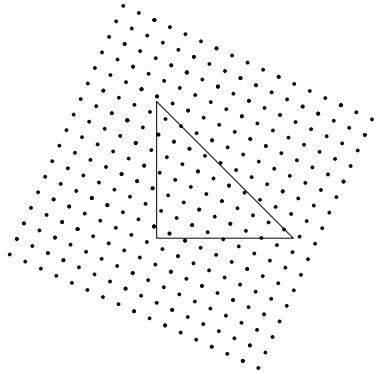




Construction Algorithm

Given a target sample size N , an angle α such as $3\pi/8$ satisfying Lemma 3. and a target triangle $\Delta(A, B, C)$,

- Take integer grid \mathbb{Z}^2
- Rotate anti clockwise by α

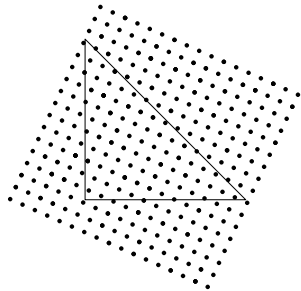




Construction Algorithm

Given a target sample size N , an angle α such as $3\pi/8$ satisfying Lemma 3. and a target triangle $\Delta(A, B, C)$,

- Take integer grid \mathbb{Z}^2
- Rotate anti clockwise by α
- Shrink by $\sqrt{2N}$

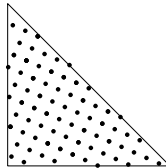




Construction Algorithm

Given a target sample size N , an angle α such as $3\pi/8$ satisfying Lemma 3. and a target triangle $\Delta(A, B, C)$,

- Take integer grid \mathbb{Z}^2
- Rotate anti clockwise by α
- Shrink by $\sqrt{2N}$
- Remove points not in the triangle.

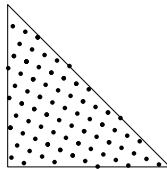




Construction Algorithm

Given a target sample size N , an angle α such as $3\pi/8$ satisfying Lemma 3. and a target triangle $\Delta(A, B, C)$,

- Take integer grid \mathbb{Z}^2
- Rotate anti clockwise by α
- Shrink by $\sqrt{2N}$
- Remove points not in the triangle.
- (Optionally) add/subtract points to get exactly N points

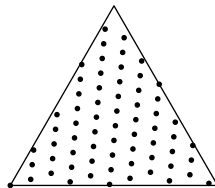




Construction Algorithm

Given a target sample size N , an angle α such as $3\pi/8$ satisfying Lemma 3. and a target triangle $\Delta(A, B, C)$,

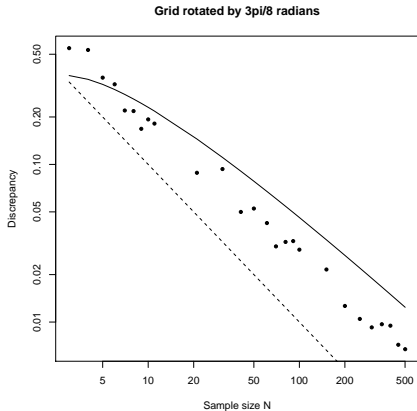
- Take integer grid \mathbb{Z}^2
- Rotate anti clockwise by α
- Shrink by $\sqrt{2N}$
- Remove points not in the triangle.
- (Optionally) add/subtract points to get exactly N points
- Linearly map R onto the desired triangle $\Delta(A, B, C)$





Triangular Kronecker Lattice

Parallel discrepancy of triangular lattice points for angle $\alpha = 3\pi/8$ and various targets N . The number of points was always N or $N + 1$. The dashed reference line is $1/N$. The solid line is $\log(N)/N$.





Conclusion

- The Kronecker construction attains a lower discrepancy than the van der Corput construction.
- van der Corput construction is extensible and the digits in it can be randomized.
- If f is continuously differentiable, then for $N = 4^k$, the randomization in Owen (1995) will give root mean square error $O(1/N)$

Future Work

- Generalization to higher dimensional simplex.
- Construction in tensor product spaces.

Thank you. Questions?